Lect. (4) 9/3/2015

⇒ Fractional numbers

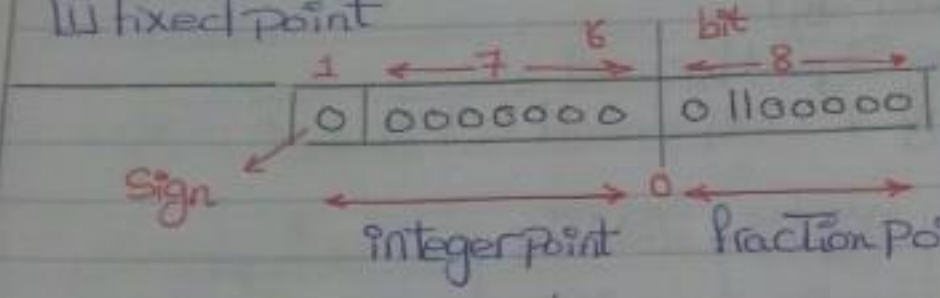
- fixed point
- floating point

⇒ Flow Chart

⇒ G - Programming

- In Computer we represent fractional numbers in 2 Ways

1) Fixed point



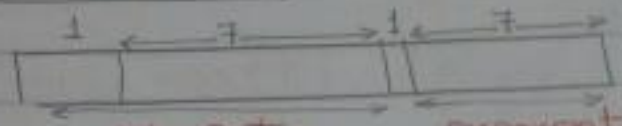
- This is the easiest but not flexible for very small and very large values.

Sign bit: 0 → +ve & 1 → -ve

⇒ 3/8 in binary

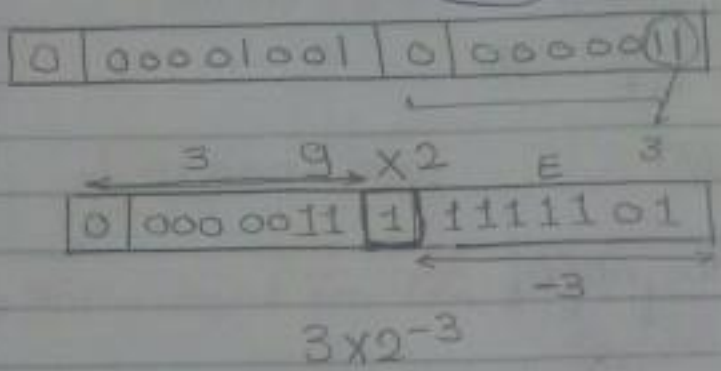
- maximum ⇒ 1111 111 - 1111 1111 ≃ 128
- minimum ⇒ 0000 000 - 0000 0001 (+ve) number (2^{-8})

[2] Floating point :-



$M * 2^E$

	Sign	
	M	E
-0.001	-ve	-ve
0.001	+ve	-ve
1000	+ve	+ve
-1000	-ve	+ve



⇒ Find 2's Complement of (3)?

3 | 2 1
1 | 2 1
0 |

8bits

0000 0011

2methods

1111 1100 → 1+ → 1111 1101

8bits

n7ot (1) ⇒ mn b3dh n8yr
el 0 → 1 & 1 → 0

1111 1101

* In C&
float = 16 bits
double = 32 bits

Maximum → |27| → 127×2^{127}

Minimum → 1×2^{-128}

⇒ $M * 2^E$ » 7bits

M
↓
7bits

0 111 1111	0 111 1111
↓	↓
127	127

Max. = 127×2^{127}

0 000 0001	1000 0000
↓	↓
1	-128

Min. = 1×2^{-128}

⇒ Add: $3 \times 2^5 + 7 \times 2^4$ → nsle Power el(2) wa7d. 3ashan n2dr n3m l add

M	x2	E
0 000 0011		0 0000 101 → Power = 5
M 0110		E 1000
0 0000 111		0 00 0010 0

$3 \times 2^5 = 3 \times 2^4 \times 2$

$= 6 \times 2^4$

0 000 0110	0 000 0100
+	
0 000 0111	0 000 0100
0 000 1101	0 000 0100

another Solution

Hasitha Power(5)

$7 \times 2^4 \rightarrow (3) \times 2^5$

0000 0011	0000 0101
000 000 11	0000 0101
0000 0110	0000 0101
0 | 0001 1000	0000 0100

6×2^5

⇒ Normalized Fraction :

Floating point

M	E
0	

must be 0

(1)

* $3/8 \Rightarrow (0.011)_B \times 2^0$

0 | 0110 0000	0 | 000 0000

Normalized (M)

0 | 1100 000	1 11 11 111